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Batching and sequencing in Bucket Brigade systems with heterogeneous items

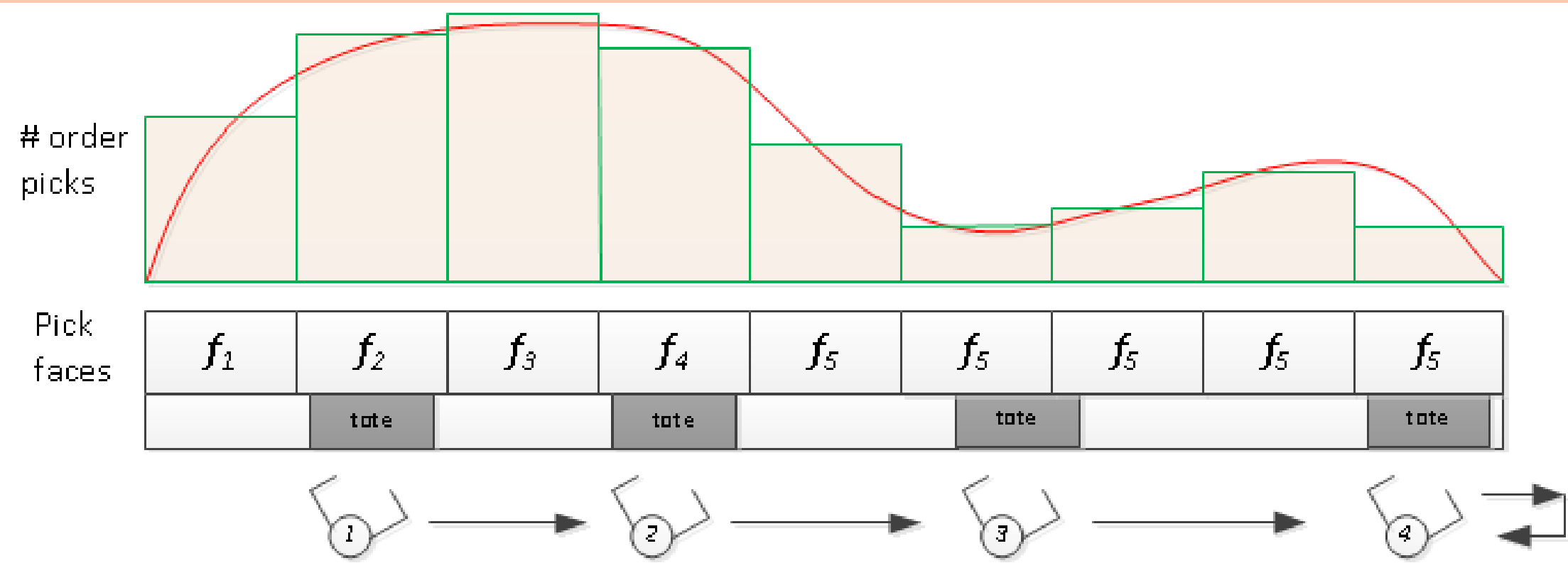
Yossi Bukchin, Eran Hanany and Eugene Khmelnitsky
Tel-Aviv University



Objectives

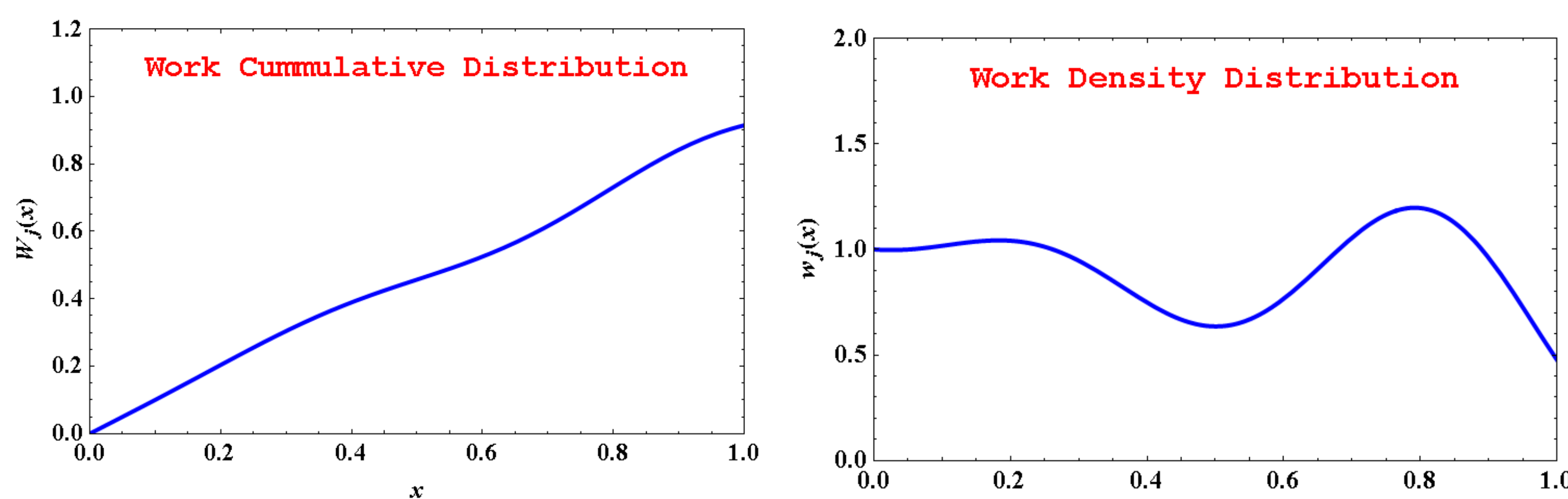
- Analyze analytically no-blockage conditions for items with different work distributions
- Accordingly, suggest order batching & sequencing approach

Heterogeneous items/orders



Analytical model: description & assumptions

- 2-worker Bucket Brigade line
- Slowest to fastest order, v_i the deterministic speed of worker i
- Multiple items to be produced /orders to be picked
- Each item/order has heterogeneous work distribution

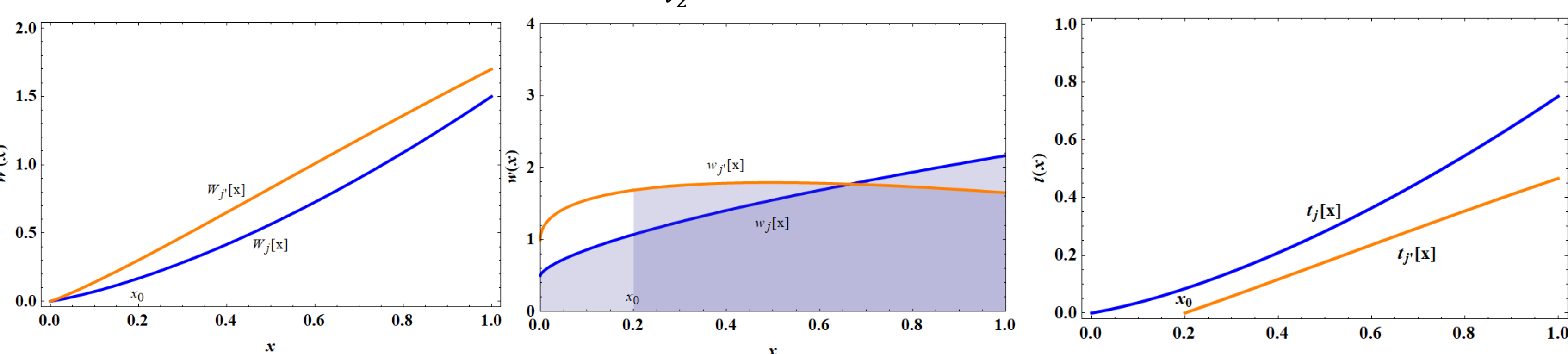


General conditions for no blockage

- $W_j[x]$ cumulative work dist. of current item j
- $W_{j'}[x]$ cumulative work dist. of new item j'
- $w_j[x] = W_j'[x]$ density work distribution
- x_0 hand-off point b/w j and j'

No-blockage condition: $\frac{W_j[x] - W_j[x_0]}{v_2} \leq \frac{W_{j'}[x]}{v_1} \quad \forall x \in [x_0, 1]$

Equivalently: define $g_{j,j'}[x] := W_j[x] - \frac{W_{j'}[x]}{v_1}$. Then, $\max_{x_0 < x \leq 1} g_{j,j'}[x] \leq W_j[x_0]$.



Sequence with no blockage when the total workload is identical

A single item type:

given a set of items, each with a work distribution $W[x]$, and $W[1] = a$, no blockage occurs. The hand-off point converges to the steady state point, x^* , which satisfies $\frac{W[x]}{v_1} = \frac{a - W[x]}{v_2}$.

$x^* = W^{-1} \left[\frac{av_1}{v_1 + v_2} \right]$ (generalization of Bartholdi Eisenstein, 1996).

Multiple item types, identical workload $a_j = a \quad \forall j$:

The steady state hand-off point of each item j , x_j^* , satisfies $\frac{W_j[x_j]}{v_1} = \frac{a - W_j[x_j]}{v_2}$.

$x_j^* = W_j^{-1} \left[\frac{av_1}{v_1 + v_2} \right]$.

no blockage occurs when the items are ordered in a sequence with weakly decreasing x_j^* .

The process starts in $\max_j x_j^*$, and each hand-off will occur in x_j^* .

Sets of item types satisfying universal no-blockage

Definition: a set of item types satisfies *universal no-blockage* if no blockage occurs for any starting position, any pair of items from the set, and any order of that pair.

For some distribution function $W^0[x]$, the set of all distribution functions $W[x]$, such that

$W[x] \in \left[\frac{v_1}{v_2} W^0[x], W^0[x] \right]$ for all $x \in [0,1]$ satisfies universal no-blockage.

Proof: $\max_{x_0 \leq x \leq 1} g_{j,j'}[x] \leq \max_{x_0 \leq x \leq 1} \left(W^0[x] - \frac{v_1 W^0[x]}{v_2} \right) = 0 \leq W_j[x_0]$.

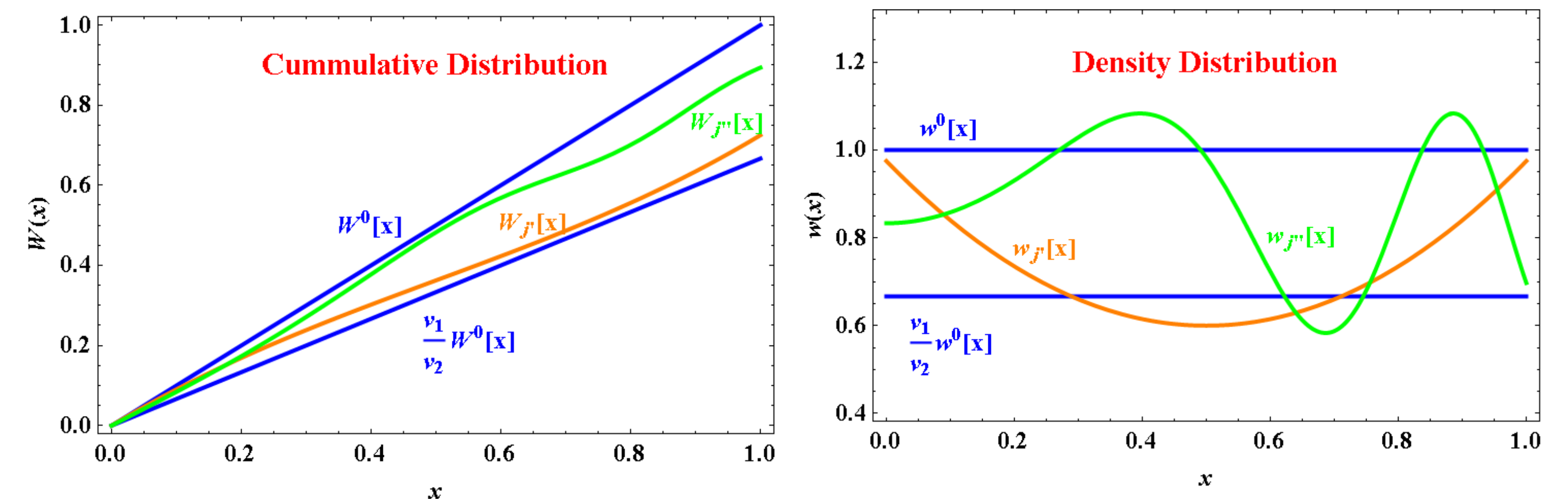
Maximal set within a domain with universal no blockage

Definition: given a domain (a set of item types), a set of item types satisfies *maximal universal no-blockage within the domain* if:

- It satisfies universal no-blockage
- For any distribution within the domain and outside the set, there exists a starting position, an item with work distribution within the set, and an order of these two items that generate blockage.

Domain of general distributions

For any distribution function W^0 , the set of all distribution functions $W[x] \in \left[\frac{v_1}{v_2} W^0[x], W^0[x] \right]$ for all $x \in [0,1]$ satisfies maximal universal no-blockage within the domain of general distributions.



Order batching

Constraint Programming for order batching

$$\text{minimize } \sqrt{\sum_{k=1}^{UB^b} \sum_{j=1}^m z_k (z_{kj} - c_j^b)^2}$$

s.t.

$$z_{kj} = \sum_{i=1}^n (x_i = k) c_{ij} \quad k = 1..UB^b, j = 1..m$$

$$\sum_{j=1}^m z_{kj} \leq C \quad k = 1..UB^b$$

$$z_k \leq \sum_{i=1}^n (x_i = k) \quad k = 1..UB^b$$

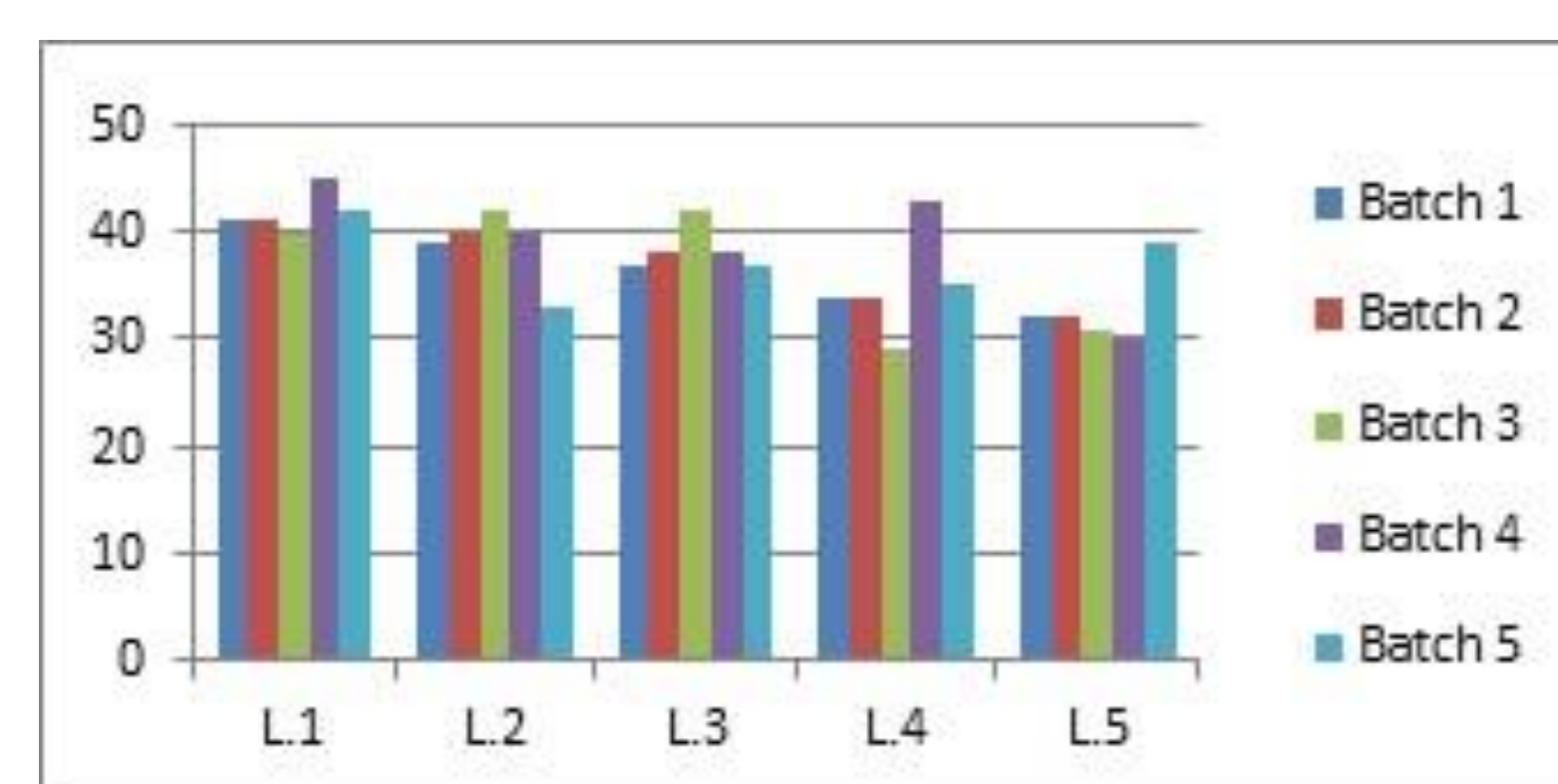
$$z_k \geq (x_i = k) \quad i = 1..n, k = 1..UB^b$$

$$z_k \geq z_{k+1} \quad k = 1..UB^b - 1$$

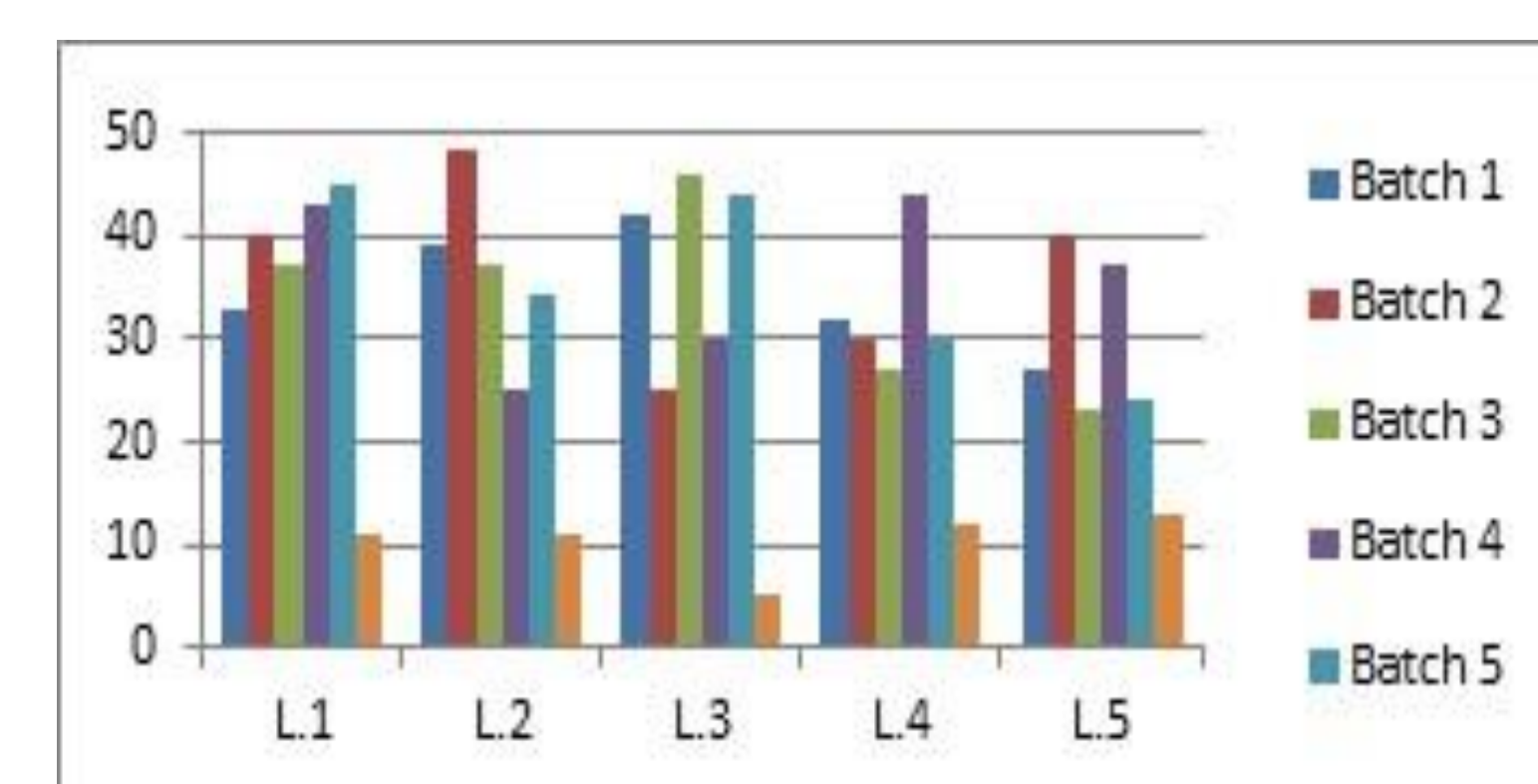
$$x_i = (0, UB^b)$$

$$z_{kj} = \text{int.}$$

$$z_k = (0,1)$$



CP batching solution (C=200)



Arbitrary batching solution (C=200)

Hamiltonian path for order sequencing